

Kenmerk: EWI2018/TW/DMMP/MU/Mod7/Exam1

**Exam 1, Module 7, Code 201700304**  
**Discrete Structures & Efficient Algorithms**  
Monday, March 19, 2018, 08:45 - 11:45

All answers need to be motivated. No calculators. You are allowed to use a handwritten cheat sheet (A4, both sides) per topic (ADS, DM, L&M).

This exam consists of three parts, with the following (estimated) times per part:

Algorithms & Data Structures (ADS)	1h	(30 points)
Discrete Mathematics (DW)	1h 20 min	(40 points)
Languages & Machines (L&M)	40 min	(20 points)

The total is  $30+40+20=90$  points. Your exam grade is the maximum of 1 and the total number of points divided by 9, rounded to one digit.

**Important:** It is necessary to use a new sheet of paper for each part (ADS/DW/L&M)!

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## Algorithms & Data Structures

- (10 points) Consider the following algorithm (where  $*$  is multiplication,  $//$  integer division (eg.  $7 // 2 = 3$ ), and  $**$  is exponentiation:

```
def func(n):
    if n==0:
        return 1
    else:
        if n<4:
            return n
        else:
            return 3*func(n//4) + 6 + func(n//4)**2
```

- Give a recursive expression for the time complexity  $T(n)$  of this algorithm (measured in the number of arithmetic operations).
  - Use the recursive expression for determining the asymptotic complexity.
- (5 points) Given a maxheap  $E$  with  $n$  elements. Give an algorithm that returns the difference between the maximum and the minimum in this heap. The algorithm should make no more than  $n/2$  comparisons.
  - (15 points) A game is played in a garden divided into  $n$  times  $n$  squares. In each square there is a number of pearls (at least 1 per square). You start in the leftmost lowest square (that has coordinates  $(1,1)$ ) and you are allowed each time to move upward or to the right. You have to end in the rightmost highest square (that has coordinates  $(n,n)$ ). You may take the pearls in each square that you pass. The goal is to get as much pearls as possible.

- (a) Let  $c(i, j)$  be the number of pearls in square  $(i, j)$ , and  $P(i, j)$  be the optimal total number of pearls you have collected when you arrive in square  $(i, j)$ . Assume  $P(i, j) = 0$  for  $i = 0$  or  $j = 0$ . Motivate which of the following recurrence relations hold for  $1 \leq i \leq n$ ,  $1 \leq j \leq n$ :
- $P(i, j) = \max\{P(i-1, j + c(i, i)), P(i-1, j)\}$
  - $P(i, j) = c(i, j) + \max\{P(i-1, j), P(i, j-1)\}$
  - $P(i, j) = c(i, j) + \min\{P(i, j-1), P(i-1, j)\}$
  - $P(i, j) = \max\{P(i-1, j + c(i, j)), P(i + c(i, j), j-1)\}$
- (b) Give an algorithm to determine the maximal number of pearls you can win. The complexity may not be bigger than  $\Theta(n^2)$ .
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## Discrete Mathematics

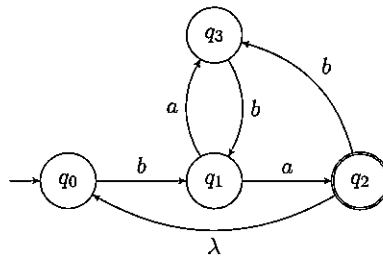
4. (5 points) Suppose we only had two different types of bills with face values  $a\text{€}$  and  $b\text{€}$ ,  $a, b \in \mathbb{Z}$ . Somebody tells you that he had recently paid  $64\text{€}$ . Another person tells that she had paid  $45\text{€}$ . Show that, for any  $k \in \mathbb{Z}$ , the amount  $k\text{€}$  can be paid using only bills of  $a\text{€}$  and  $b\text{€}$ .
5. (10 points)
- Denote by  $a_n$  the number of strings in  $\{0, 1, 2\}^*$  of length  $n$  that contain no substring  $11$ , and neither  $22$ . Compute  $a_1$  and  $a_2$ . Set up a recurrence relation for  $a_n$ ,  $n \geq 3$ . (You do not have to solve that recurrence relation.)
  - Compute the solution to the recurrence relation
 
$$a_n - 10a_{n-1} + 25a_{n-2} = 16n + 8 \quad (n \geq 2) \quad \text{with } a_0 = 3 \text{ and } a_1 = 12.$$
6. (8 points) Suppose we want to share  $100\text{€}$  among three persons, such that each of them gets at least  $20\text{€}$ , but at most  $50\text{€}$ , and moreover each person gets an integer amount. How many possibilities are there to do that? Use a generating function to compute your answer.
7. (7 points) Let  $G = (V, E)$  simple, connected, undirected graph with  $|V| = n$  and  $|E| = m$ . Show that, if  $G$  has an average node degree  $\sum_{v \in V} \frac{d(v)}{n} \geq 6$ , then  $G$  cannot be planar.
8. (10 points) Give a short proof or give a counterexample for each of the following statements.
- Consider an undirected, simple graph  $G = (V, E)$  with edge weights  $w_e \geq 0$ ,  $e \in E$ . Then any two minimum spanning trees  $T_1$  and  $T_2$  for  $G$  must have a nonempty intersection, that is,  $T_1 \cap T_2 \neq \emptyset$ .
  - Consider a capacitated network  $G = (V, A, c)$ , where  $V$  is the set of vertices,  $A$  is the set of directed arcs, and  $c_a \geq 0$ ,  $a \in A$  are the arc capacities. Then there always exists a maximum flow  $f_a$ ,  $a \in A$ , such that either  $f_a = 0$  or  $f_a = c_a$  for all  $a \in A$ .
  - Consider an undirected, simple graph  $G = (V, E)$  with edge weights  $w_e \geq 0$  such that  $w_e \neq w_{e'}$  for all  $e, e' \in E$ ,  $e \neq e'$ . Let  $s \in V$  be fixed. Then for all  $v \in V$  there is a unique shortest  $(s, v)$ -path.
  - Consider an undirected, simple graph  $G = (V, E)$  with edge weights  $w_e \geq 0$  such that  $w_e \neq w_{e'}$  for all  $e, e' \in E$ ,  $e \neq e'$ . Then there is a unique minimum spanning tree  $T$ .

- (e) In a capacitated network  $G = (V, A, c)$ , where  $V$  is the set of vertices,  $A$  is the set of directed arcs,  $c_a \geq 0$ ,  $a \in A$ , are the arc capacities, and  $c_a \neq c_{a'}$  for all  $a, a' \in A$ ,  $a \neq a'$ . Then there is a unique minimum cut.

## Languages & Machines

Please remember to start this part on a new sheet of paper!

9. (11 points) Consider the following NFA with  $\lambda$ -steps  $M$  (only  $q_2$  is accepting):



- First eliminate state  $q_3$ , adding new transitions labeled with regular expressions, to preserve the accepted language. Show the resulting "expression graph".
  - Continue the construction, by eliminating  $q_1$  as well, and read off a regular expression  $E$  with  $\mathcal{L}(E) = \mathcal{L}(M)$ .
  - Provide the  $\lambda$ -closure and input-transition function of the automaton  $M$  in a table.
  - Transform the automaton  $M$  in a systematic manner to a (possibly incomplete) DFA.
10. (9 points) Consider the definitions of the following languages over  $\Sigma = \{a, b\}$ :
- Language  $L_1 := \{a^{2i} b^j \mid 0 \leq i \text{ and } 0 \leq j\}$
  - Language  $L_2 := \{a^i b^j a^i \mid 0 \leq i \text{ and } 0 \leq j\}$
  - Language  $L_3$  is an (arbitrary) *finite* language

Indicate for each of the following languages if they are regular or not. Motivate your answers, either by a proof or a construction.

- Language  $L_1$
- Language  $L_2$
- Language  $\overline{L_3} \cup L_3^R$

