

Kenmerk : TW2018/DWMP/001/ha

Course : **Discrete Mathematics for Technical Computer Science; Part 1**

Date : October 26, 2018

Time : 08.45–09.45 hrs

**Motivate all your answers. The use of electronic devices is not allowed.
A formula sheet is included.**

In this exam: $\mathbb{N} = \{0, 1, 2, 3, \dots\}$.

1. [6 pt]

Let $f : \mathbb{R} \rightarrow \mathbb{R}$ be a function. Give quantified expressions for the following statements.

(a) f is decreasing.

(b) The minimum value of f is equal to 0 and the maximum value of f is equal to 5.

2. [6 pt]

Prove the validity of the following argument using the "Laws of Logic" and the "Rules of Inference".

$$\begin{array}{l} \neg q \vee s \\ (p \rightarrow q) \wedge r \\ \frac{(p \rightarrow r) \rightarrow p}{\therefore \neg(s \rightarrow \neg p)} \end{array}$$

3. [6 pt]

Let A , B and C be sets in a universe \mathcal{U} . Give a proof or a counterexample for each of the following statements.

(a)

$$(A - C = B - C \quad \wedge \quad A \cup C = B \cup C) \implies A = B.$$

(b)

$$(A - C = B - C \quad \wedge \quad A \cap C = B \cap C) \implies A = B.$$

Total: 18 points

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Course : **Discrete Mathematics for Technical Computer Science; Part 2**

Date : October 26, 2018

Time : 09.45–10.45 hrs

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In this exam: $\mathbb{N} = \{0, 1, 2, 3, \dots\}$.

4. [6 pt]

Let the sequence of numbers $a_0, a_1, a_2, a_3, \dots$ be given by:

$$a_0 = 1, a_1 = 1, a_2 = 2, \text{ and for } n \geq 3: a_n = a_{n-2} + a_{n-3}.$$

Prove with mathematical induction that for all $n \in \mathbb{N}$:

$$a_n \leq (\sqrt{2})^n.$$

5. [6 pt]

Let f be the closed binary operation on $\mathcal{P}(\{1, 2, 3\})$ given by

$$f(A, B) = A - B \quad (A, B \subseteq \{1, 2, 3\})$$

- (a) Determine if f is one-to-one; if f is onto; and if f is invertible.
- (b) Determine $f^{-1}(\{1, 3\})$.

6. [6 pt]

Let D be a set in a universe \mathcal{U} . Define the relation R on $\mathcal{P}(\mathcal{U})$ by:

$$ARB \text{ if and only if } A - D = B - D \quad (A, B \subseteq \mathcal{U}).$$

- (a) Show that R is an equivalence relation on $\mathcal{P}(\mathcal{U})$.
- (b) Take $\mathcal{U} = \{1, 2, 3, 4\}$ and $D = \{1, 2\}$. Determine the partition of $\mathcal{P}(\mathcal{U})$ induced by R .

Total: 18 points

Rules of Inference

R1.	$\frac{p \quad p \rightarrow q}{\therefore q}$	Modus Ponens
R2.	$\frac{p \rightarrow q \quad q \rightarrow r}{\therefore p \rightarrow r}$	Law of the Syllogism
R3.	$\frac{p \rightarrow q \quad \neg q}{\therefore \neg p}$	Modus Tollens
R4.	$\frac{p \quad q}{\therefore p \wedge q}$	Rule of Conjunction
R5.	$\frac{p \vee q \quad \neg p}{\therefore q}$	Rule of Disjunctive Syllogism
R6.	$\frac{\neg p \rightarrow F_0}{\therefore p}$	Rule of Contradiction
R7.	$\frac{p \wedge q}{\therefore p}$	Rule of Conjunctive Simplification
R8.	$\frac{p}{\therefore p \vee q}$	Rule of Disjunctive Amplification
R9.	$\frac{p \wedge q \quad p \rightarrow (q \rightarrow r)}{\therefore r}$	Rule of Conditional Proof
R10.	$\frac{p \rightarrow r \quad q \rightarrow r}{\therefore (p \vee q) \rightarrow r}$	Rule for Proof by Cases
R11.	$\frac{p \rightarrow q \quad r \rightarrow s \quad p \vee r}{\therefore (q \vee s)}$	Rule of the Constructive Dilemma
R12.	$\frac{p \rightarrow q \quad r \rightarrow s \quad \neg q \vee \neg s}{\therefore \neg p \vee \neg r}$	Rule of the Destructive Dilemma

Laws of Logic

- L1.** $\neg\neg p \Leftrightarrow p$ Law of Double Negation
- L2.** $\neg(p \vee q) \Leftrightarrow \neg p \wedge \neg q$ DeMorgan's Laws
 $\neg(p \wedge q) \Leftrightarrow \neg p \vee \neg q$
- L3.** $p \vee q \Leftrightarrow q \vee p$ Commutative Laws
 $p \wedge q \Leftrightarrow q \wedge p$
- L4.** $p \vee (q \vee r) \Leftrightarrow (p \vee q) \vee r$ Associative Laws
 $p \wedge (q \wedge r) \Leftrightarrow (p \wedge q) \wedge r$
- L5.** $p \vee (q \wedge r) \Leftrightarrow (p \vee q) \wedge (p \vee r)$ Distributive Laws
 $p \wedge (q \vee r) \Leftrightarrow (p \wedge q) \vee (p \wedge r)$
- L6.** $p \vee p \Leftrightarrow p$ Idempotent Laws
 $p \wedge p \Leftrightarrow p$
- L7.** $p \vee F_0 \Leftrightarrow p$ Identity Laws
 $p \wedge T_0 \Leftrightarrow p$
- L8.** $p \vee \neg p \Leftrightarrow T_0$ Inverse Laws
 $p \wedge \neg p \Leftrightarrow F_0$
- L9.** $p \vee T_0 \Leftrightarrow T_0$ Domination Laws
 $p \wedge F_0 \Leftrightarrow F_0$
- L10.** $p \vee (p \wedge q) \Leftrightarrow p$ Absorption Laws
 $p \wedge (p \vee q) \Leftrightarrow p$
- L11.** $p \leftrightarrow q \Leftrightarrow (p \rightarrow q) \wedge (q \rightarrow p)$
- L12.** $p \rightarrow q \Leftrightarrow \neg p \vee q$
- L13.** $p \rightarrow q \Leftrightarrow \neg q \rightarrow \neg p$

Laws concerning quantifiers

$$\text{N1. } \neg[\forall x p(x)] \Leftrightarrow \exists x \neg p(x)$$

$$\text{N2. } \neg[\exists x p(x)] \Leftrightarrow \forall x \neg p(x)$$

Additional Laws concerning quantifiers

$$\text{U1. } \frac{\forall x p(x)}{\therefore p(c) \text{ for arbitrary } c \text{ in the universe}}$$

$$\text{U2. } \frac{\exists x p(x)}{\therefore p(c) \text{ for some } c \text{ in the universe}}$$

$$\text{U3. } \frac{p(c) \text{ for arbitrary } c \text{ in the universe}}{\therefore \forall x p(x)}$$

$$\text{U4. } \frac{p(c) \text{ for some } c \text{ in the universe}}{\therefore \exists x p(x)}$$

U1: Rule of Universal Specification

U2: Rule of Existential Specification

U3: Rule of Universal Generalization

U4: Rule of Existential Generalization