

Examination Secure Data Management

192110940 (UT students)

192110941 (Kerckhoff students)

- November 1st, 2013 -

Instruction:

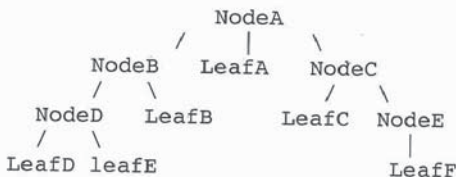
- This is an open book examination
- No electronic devices allowed
- The examination consists of TWENTY multiple-choice questions; each has the same weight: 5 POINTS
- Success!

1/ Assume that all devices have the same probability of being revoked. Which device in the following media key block has the highest chance of getting blocked due to **other** devices being revoked?

WAL GEP ADE
 CPL CAP WDC
 DAL BAC GPE

- A/ D
- B/ G
- C/ P
- D/ W

2/ What is the Prime p of the Finite Ring $\mathbb{F}_p[x]/(x^{p-1} - 1)$ used for searching in the following XML structure?



- A/ $p = 7$
- B/ $p = 11$
- C/ $p = 13$
- D/ $p = 17$

3/ Which version of OMA supports fully secure super distribution?

- A/ OMA v1 Combined Delivery mode
- B/ OMA v1 Separate Delivery mode
- C/ OMA v2 Group Functionality
- D/ None of the above

4/ What is the key difference between Copy Protection and DRM?

- A/ Copy Protection rules Copy Control while DRM rules Usage Control
- B/ Copy Protection is used for discs; DRM is used for Internet
- C/ Copy Protection rules Copy Control while DRM rules Access Control
- D/ There is no fundamental difference

5/ Assume attributes with corresponding terms A - t1, B - t2, C - t3, D - t4, E - t5 and F - t6. Which Subjects will get access under the policy encoded as $c = (\tau, g^{12}, m * e(g, g)^{12 * \alpha} g^{5 * t1}, g^{3 * t2}, g^{3 * t3}, g^{2 * t4}, g^{2 * t5}, g^{5 * t6})$?

- A/ Those that have at least attributes A, B and E
- B/ Those that have at least attributes F, C, D and E
- C/ Those that have at least attributes B, D and E
- D/ Those that have at least attributes A and D

6/ Given the following SQL query "SELECT sub1, sub2 FROM tab1, tab2 WHERE sub1.att=sub2.att". Assume the following mapping functions.

tab1.att	tab2.att
[100 - 200] -> 1	[100 - 250] -> 9
[200 - 300] -> 3	[250 - 500] -> 6
[300 - 400] -> 8	
[400 - 500] -> 7	

What is the correct representation of the query condition on the encrypted tables?

- A/ $((sub1)^s.att=1 \text{ AND } (sub2)^s.att=9) \text{ OR } ((sub1)^s.att=3 \text{ AND } (sub2)^s.att=9) \text{ OR } ((sub1)^s.att=8 \text{ AND } (sub2)^s.att=6) \text{ OR } ((sub1)^s.att=7 \text{ AND } (sub2)^s.att=6)$
- B/ $(100 < (sub1)^s.att < 200 \text{ AND } 100 < (sub2)^s.att < 250) \text{ OR } (200 < (sub1)^s.att < 300 \text{ AND } 100 < (sub2)^s.att < 250) \text{ OR } (200 < (sub1)^s.att < 300 \text{ AND } 250 < (sub2)^s.att < 500) \text{ OR } (300 < (sub1)^s.att < 400 \text{ AND } 250 < (sub2)^s.att < 500) \text{ OR } (400 < (sub1)^s.att < 500 \text{ AND } 250 < (sub2)^s.att < 500)$
- C/ $((sub1)^s.att=1 \text{ AND } (sub2)^s.att=9) \text{ OR } ((sub1)^s.att=3 \text{ AND } (sub2)^s.att=9) \text{ OR } ((sub1)^s.att=3 \text{ AND } (sub2)^s.att=6) \text{ OR } ((sub1)^s.att=8 \text{ AND } (sub2)^s.att=6) \text{ OR } ((sub1)^s.att=7 \text{ AND } (sub2)^s.att=6)$
- D/ $((100 <= (sub1)^s.att < 200 \text{ AND } 100 <= (sub2)^s.att < 250) \text{ OR } (200 <= (sub1)^s.att < 300 \text{ AND } 100 <= (sub2)^s.att < 250) \text{ OR } (200 <= (sub1)^s.att < 300 \text{ AND } 250 <= (sub2)^s.att < 500) \text{ OR } (300 <= (sub1)^s.att < 400 \text{ AND } 250 <= (sub2)^s.att <= 500) \text{ OR } (400 <= (sub1)^s.att <= 500 \text{ AND } 250 <= (sub2)^s.att <= 500)$

7/ What is the effect of removing the random value r in the ABE access control term $d_0 = g^{-r}$?

- A/ Removing r will block access to attribute r
- B/ Removing r allows users to collude
- C/ Removing r will make the scheme more efficient
- D/ Removing r reduces the number of policies that are supported

8/ Which role does the DDH assumption play in game based security analysis of the El Gamal scheme?

- A/ It serves to prove that $\Pr[\text{Attacker succeeds in game 1}] - \Pr[\text{Attacker succeeds in game 2}]$ is negligible
- B/ It serves to prove that $\Pr[B(g^x, g^y, g^{xy}) = 1] = 1/2 + \epsilon$
- C/ It serves to prove that $|\Pr[B(g^{z^1}, g^{z^2}, g^{z^1z^2}) = 1] - \Pr[B(g^{z^1}, g^{z^2}, Z) = 1]|$ is negligible
- D/ It serves to prove that $Enc_{mi} = (g^x, m_i g^{xy})$

9/ Which point below is **NOT** on the (3,6) Shamir Secret Sharing polynomial with secret 67 and random $a_1=3$ and $a_2=7$? \leftarrow faut

- A/ (2, 65)
- B/ (7, 165)
- C/ (6, 131)
- D/ (5, 107)

$$a_1 = -7 \quad a_2 = 3$$

10/ In the wildcard search what will be the effect on the term below in case there are no wildcards in the query word?

$$\prod_{\substack{i=1 \\ i \notin J}}^n U_i^{w_i' \prod_{j \in J} (i-j)}$$

- A/ The term will evaluate to 0
- B/ The term will evaluate to 1
- C/ The term cannot be evaluated
- D/ The term will evaluate to a specific number representing the query word

11/ Delegated search: how will the Delegate step $t^*=(t_1, t_2, t_3, t_4)=$
 $(\gamma^a \cdot pk_S^{r1}, \gamma^{r1}, \gamma^{y^a} \cdot pk_S^{r2}, \gamma^{r2})$ look like in case delegation is based on group G?

A/ $t^*=(t_1, t_2, t_3, t_4)= (\gamma^a \cdot pk_g^{r1}, \gamma^{r1}, \gamma^{y^a} \cdot pk_g^{r2}, \gamma^{r2})$

B/ $t^*=(t_1, t_2, t_3, t_4)= (g^a \cdot pk_S^{r1}, g^{r1}, g^{y^a} \cdot pk_S^{r2}, g^{r2})$

C/ $t^*=(t_1, t_2, t_3, t_4)= (\gamma^a \cdot pk_S^g, \gamma^g, \gamma^{y^a} \cdot pk_S^{r2}, \gamma^{r2})$

D/ $t^*=(t_1, t_2, t_3, t_4)= (g^b \cdot pk_S^{r1}, g^{r1}, g^{g^b} \cdot pk_S^{r2}, g^{r2})$

12/ Delegated search: which step in the scheme would not work properly in case the symmetry property of the bi-linear mapping does not hold?

- A/ KeyGen
- B/ Delegate
- C/ TrapGen
- D/ Test1

13/ Which of the following statements is **FALSE**?

- A/ Bilinear maps exist on certain elliptic curves
- B/ A commitment scheme cannot be both computationally hiding and computationally binding
- C/ In an unconditionally binding commitment scheme, no computationally limited adversary can reveal two different values
- D/ The shares in the Blackley secret sharing scheme are hyper planes

14/ What is the role of a ROM mark?

- A/ A ROM mark prevents bit-wise copies because the different ROM mark of the destination disc prevents proper decryption
- B/ A ROM mark prevents copying a disc bit-wise because the ROM mark blocks copying
- C/ A ROM mark prevents bit-wise copies because the ROM mark of the original disc prevents any decryption
- D/ A ROM mark prevents copying a disc bit-wise due to the missing ROM mark of the destination disc preventing encryption

15/ Which of the following functions over the natural numbers is a homomorphism from addition to multiplication?

- A/ $f(x) = x^2$
- B/ $f(x) = 2x$
- C/ $f(x) = 2^x$
- D/ $f(x) = x+2$

16/ Given a path $p = \text{label}_1/\text{label}_2/\dots/\text{label}_5$, and a function $g(x)$ which turns a label into a natural number representing the binary representation of the first character of the label. Which hash function $h(p)$ is expected to have least collisions?

- A/ $h(p) = (f(\text{label}_1)*f(\text{label}_2)*\dots*f(\text{label}_5)) \text{ MOD } 13$, where $f(x) = 2*g(x)$
- B/ $h(p) = (f(\text{label}_1)*f(\text{label}_2)*\dots*f(\text{label}_5)) \text{ MOD } 11$, where $f(x) = g(x) \text{ MOD } 2$
- C/ $h(p) = (f(\text{label}_1)*f(\text{label}_2)*\dots*f(\text{label}_5)) \text{ MOD } 11$, where $f(x) = (g(x))^2$
- D/ $h(p) = (f(\text{label}_1)*f(\text{label}_2)*\dots*f(\text{label}_5)) \text{ MOD } 13$, where $f(x) = g(x) \text{ MOD } 2$

17/ Which of the following statements is false?

- A/ Mandatory Access Control (MAC) is used in most modern consumer operating systems.
- B/ Role Based Access Control (RBAC) is particularly suitable for a company with a high turnover.
- C/ In biometrics, false-positives are much worse than false-negatives.
- D/ Biometric properties can change over time.

18/ Consider the Search in Encrypted Data approach from Song, Wagner, Perrig. Which statement below is **FALSE**?

- A/ The search is linear.
- B/ The approach exploits $X \text{ xor } K = C$ and $C \text{ xor } X = K$.
- C/ The approach exploits the homomorphic property of K .
- D/ The approach exploits $X = X1 | X2$ and $X1 = F(X2)$.

19/ An RSA SecurID token is an example of an authentication method that is:

- A/ Asynchronous.
- B/ Based on a shared secret between the server and a token.
- C/ Both A/ and B/.
- D/ Neither A/ or B/.

20/ Which of the following sets forms a CYCLIC GROUP with the given algebraic operation?

- A/ Z_n^* with addition MOD n for $n \in N$
- B/ Z_{15}^* with multiplication MOD 15
- C/ Z_6^* with multiplication MOD 6
- D/ Z_p with multiplication MOD p for a prime p

Toets Parel 000 der Informatica (201300070)

6 september 2013, 10:45–11:45

- Je mag 1 zelfgemaakt A4'tje met aantekeningen bij dit tentamen gebruiken. Rekenmachines, laptops, mobiele telefoons e.d. zijn niet toegestaan. **Stop deze nu in je tas!**
- Het aantal te behalen punten per opgave staat in de marge.

1. Binaire getallen

- (a) Reken het 1-complements binaire getal 10101 om naar decimaal. Laat zien hoe je dit berekent. 7
- (b) Reken het hexadecimale getal A5 om naar decimaal. Laat zien hoe je dit berekent. 7
- (c) Stel je hebt een 8-bits unsigned binair getal en je schuift alle bits 1 plek naar rechts, en vult links een 0 aan. Met welke rekenkundige bewerking correspondeert dit, en waarom? 7
- (d) Verder over de vorige deelvraag: werkt dit ook goed als het 8-bits getal een signed 2-complements getal zou zijn? Waarom? 4

2. Booleaanse logica

Iemand wil een schakeling bouwen die van twee input-bits A en B, naar keuze de OR of de AND berekent, onder besturing van een derde inputbit C; C=0 voor OR en C=1 voor AND.

- (a) Geef de waarheidstabel hiervan. 8
- (b) Vereenvoudig de volgende Booleaanse formule en geef daarbij aan welke Booleaanse rekenregel(s) je gebruikt; begin met het wegwerken van het '+'-teken.
 $A + A \cdot \bar{B}$ 9
- (c) Schets hoe je met alleen NOR-poorten de volgende formule kunt realiseren:
 $A + B + \bar{C}$ 8

Zie volgende bladzijde...

3. Opgave 3

20



Bovenstaande simpele processor kent twee instructies: 0 = '+' (optellen) en 1 = '*' (vermenigvuldigen).

Geef voor deze processor het programma voor de volgende bewerking: $R2 = (R1+R2)*(R0+R1)$

	leesadres 1 / schrijfadres	leesadres 2	instructie
Tijdslot 0			
Tijdslot 1			
Tijdslot 2			
Tijdslot 3			
Tijdslot 4			
...			

4. Opgave 4

30

Gegeven het volgende AVR-programma ("BRNE" betekent "BRanch if Not Equal", "MUL" betekent MULTIply, "DEC" betekent "DECrement (verminderen met 1)" en "SUB" betekent "Subtract"):

```
LDI R16, $03
LDI R17, $03
LDI R18, $02
LDI R20, $01
MUL R17, R18
DEC R16
MOV R19, R16
SUB R19, R20
BRNE -5
```

Geef in tabelvorm aan hoe de inhoud van de registers verandert terwijl dit programma doorlopen wordt; dat mag hexadecimaal of decimaal, naar keuze, maar laat wel zien wat je kiest.

Einde van deze toets.